Deniability in Automated Contact Tracing

Christoph Günther

Krzysztof Pietrzak

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Automated contact tracing (ACT)

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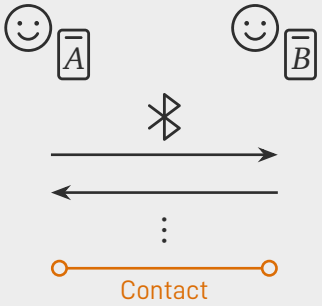
- A lot of research & development due to Covid-19
- Google and Apple jointly deployed protocol based on DP3T (Troncoso et al.)

ACT using Bluetooth Low Energy



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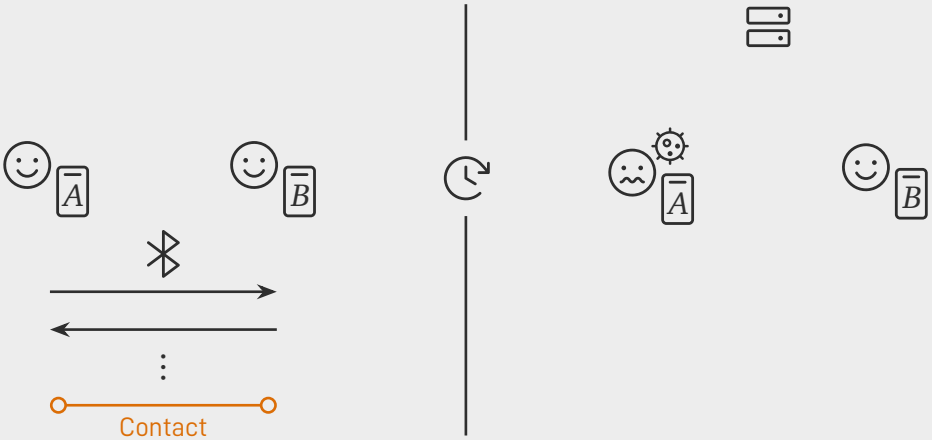

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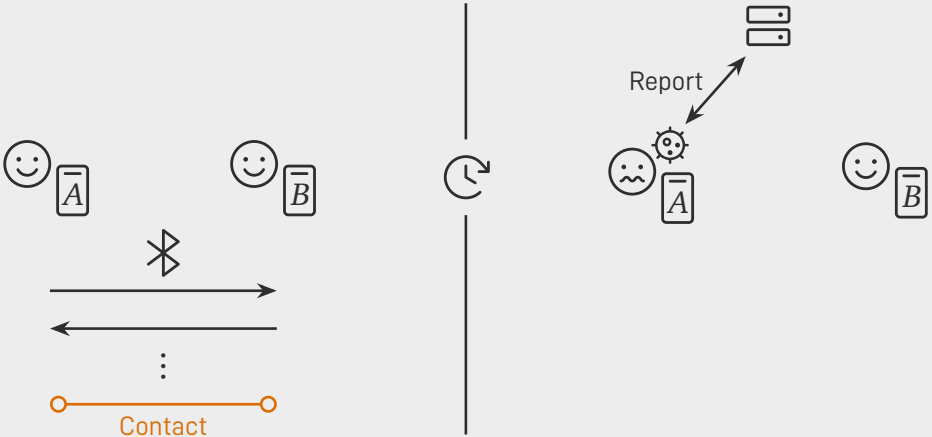

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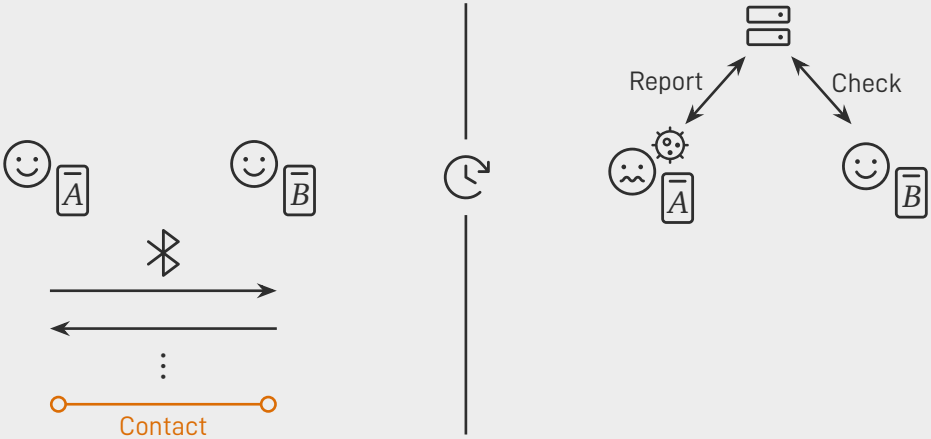

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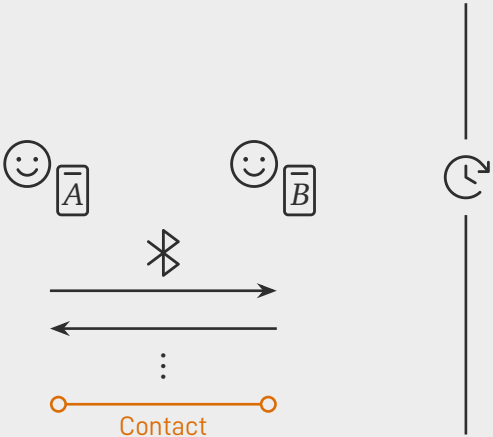
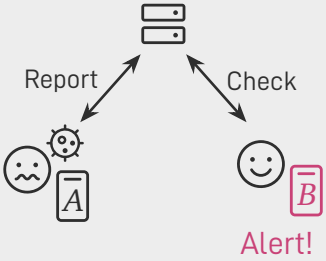


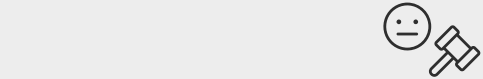


Contact and A reported sick
$\Rightarrow$
B's check causes an alert

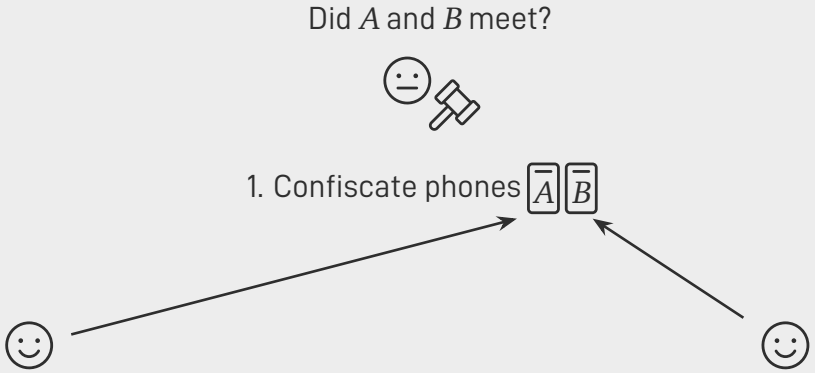
Privacy?

Digital evidence

Did A and B meet?



Digital evidence



Digital evidence

Did A and B meet?

1. Confiscate phones $\boxed{\overline{A}}$ $\boxed{\overline{B}}$
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$\boxed{\overline{B}}$ got an alert $\Leftrightarrow$ A and B have met

Deniability

- User cannot deny having met by definition of ACT

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Contact-time deniability

Did A and B meet at time t?

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Contact-time deniability

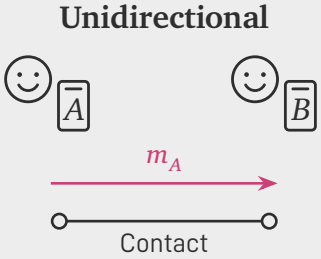
Did A and B meet at time t?

Impossibility result

Dilemma between contact-time deniability and security for a large class of protocols capturing many practically relevant ones

Impossibility result

Practical protocols are often…



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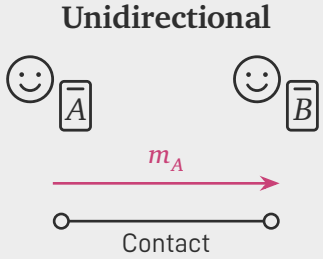

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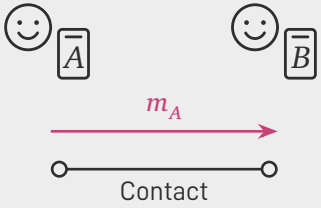
Practical protocols are often…

Unidirectional



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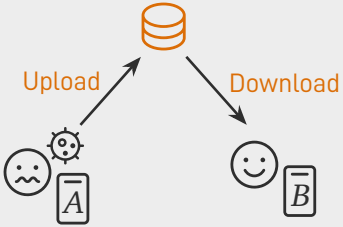
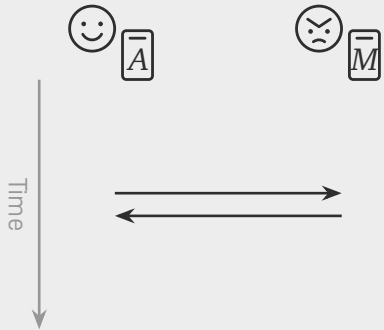
Decentralized



Server is a database

Contact-time deniability



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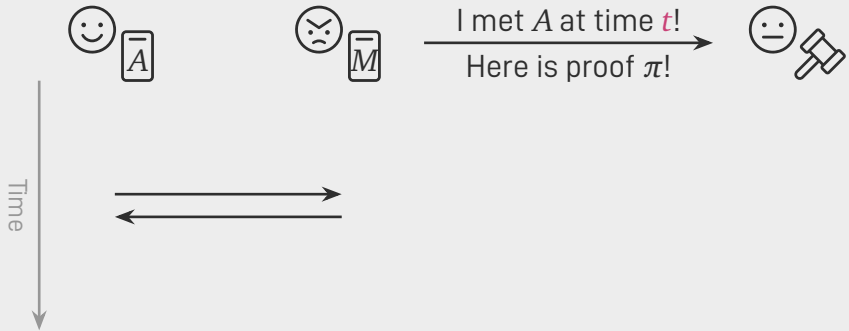

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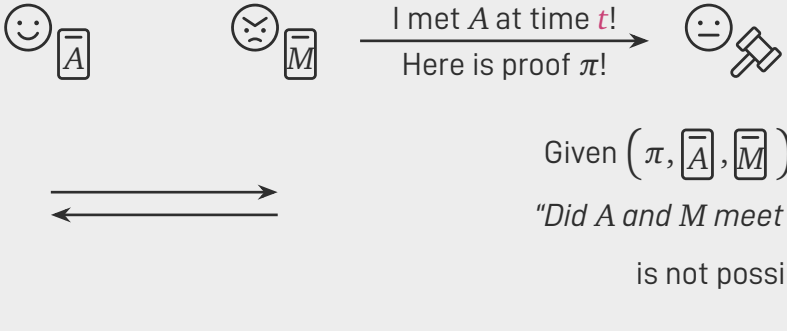


Given $\left(\pi, \bar{A}, \bar{M}\right)$, deciding

"Did A and M meet at time t?"

is not possible

Δ-contact-time deniability

Given $\left(\pi, \bar{A}, \bar{M}\right)$, deciding

"Did A and M meet at time $t \pm \Delta$*?"*

is not possible

Δ-contact-time deniability

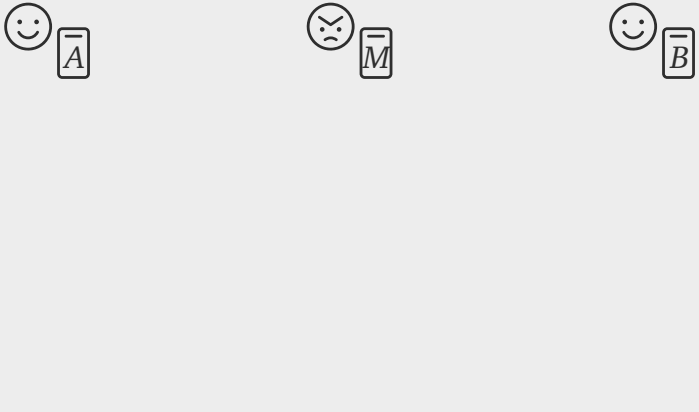
Given $\left(\pi, \overline{A}, \overline{M}\right)$, deciding

"Did A and M meet at time $t \pm \Delta$*?"*

is not possible

Meaningful if large (e.g., a week)

Δ-replay security



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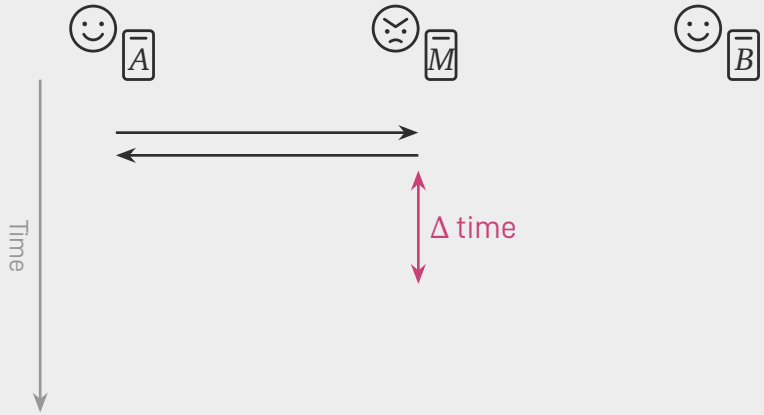

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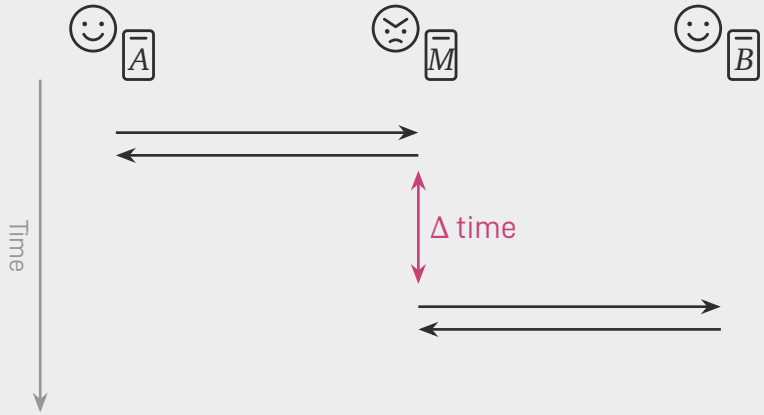

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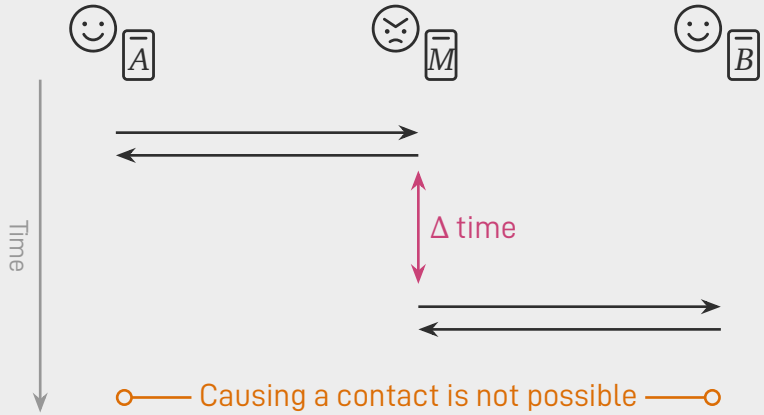

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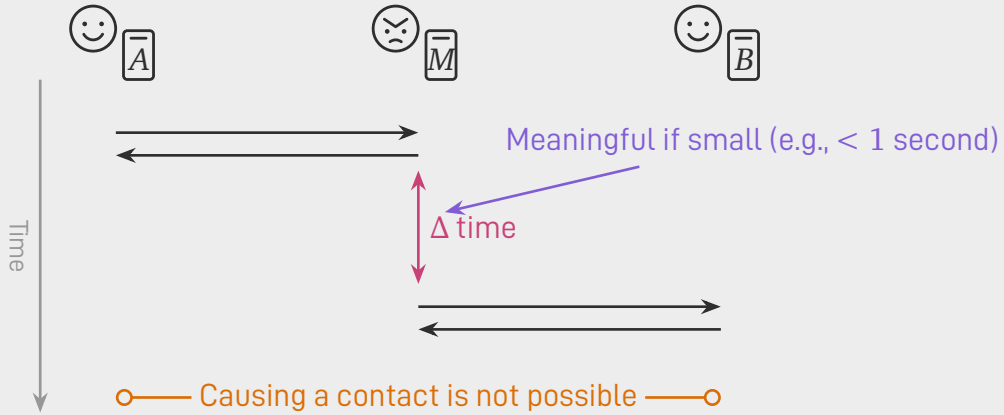

Main theorem

If a unidirectional, decentralized protocol is Δ-replay secure, then it can at most be Δ'-contact-time deniable with $\Delta > \Delta'$.

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If a unidirectional, decentralized protocol is Δ-replay secure, then it can at most be Δ'-contact-time deniable with $\Delta > \Delta'$.

Problematic since

- Δ only meaningful if small
- Δ' only meaningful if large

Proof

Want to show: Δ-replay secure ⇒ **not** Δ-contact-time deniable

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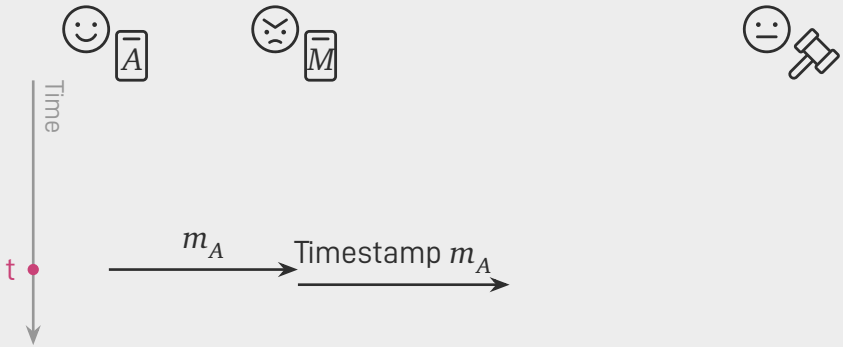

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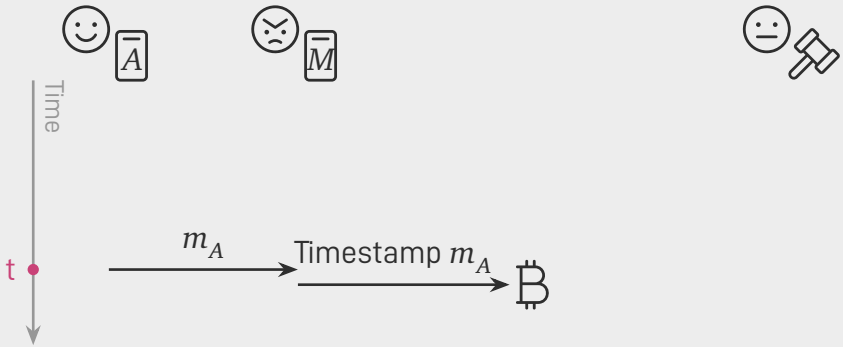

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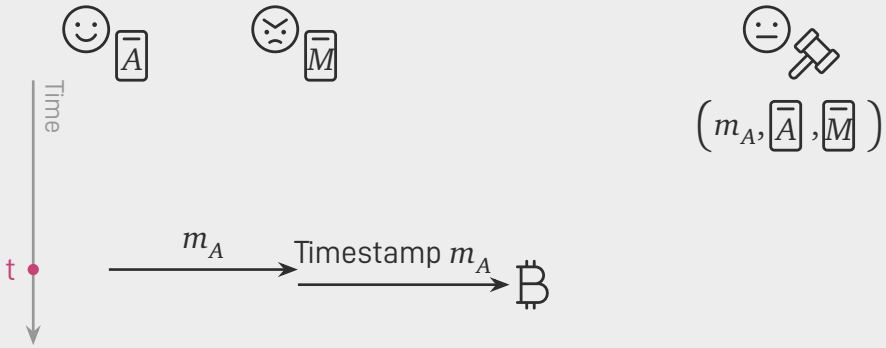

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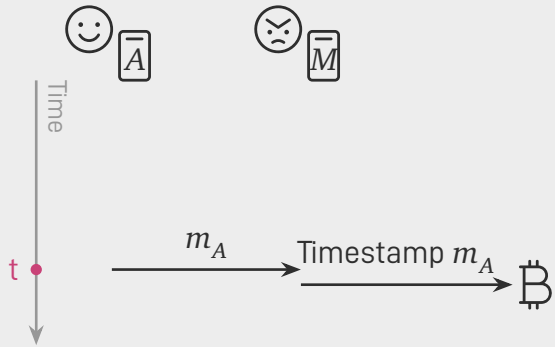
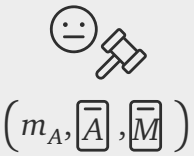
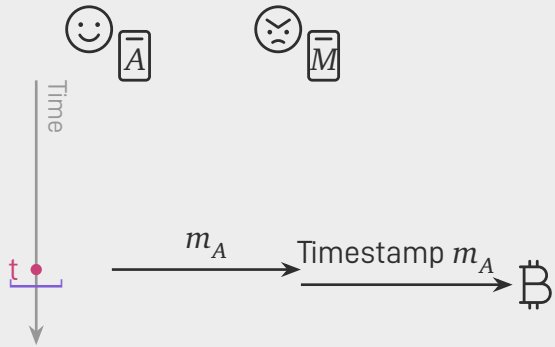




1. Would receiving m_A at time *t* lead to contact with $\bar{A}$?

Proof

Want to show: Δ-replay secure ⇒ **not** Δ-contact-time deniable



$(m_A, \bar{A}, \bar{M})$

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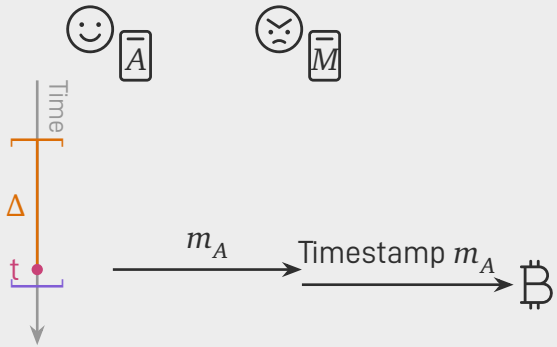

$(m_A, \bar{A}, \bar{M})$

1. Would receiving m_A at time t lead to contact with $\bar{A}$?

2. $m_A \in$ ₿ at time t?

3. Δ-replay attack not possible!

Various combinations exist

Protocol	Unidirectional	Replay-secure	Deniable
DP3T	✓	∼	∼
Challenge-Resp.	×	✓	✓
Delayed Auth.	✓	✓	×
CleverParrot	✓	✓	×
NTK	✓	∼	∼
DH-based	×	✓	✓

Table: Properties of existing decentralized protocols (cf. paper for details).

Conclusion

- Deniability not achievable for large part of the design space
- Extends to all decentralized ACT assuming stronger judge (cf. paper)
- How does this help in practice? (cf. paper)
  - Identify interesting points in the design space
  - Find creative ways to break out the theoretical model

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